

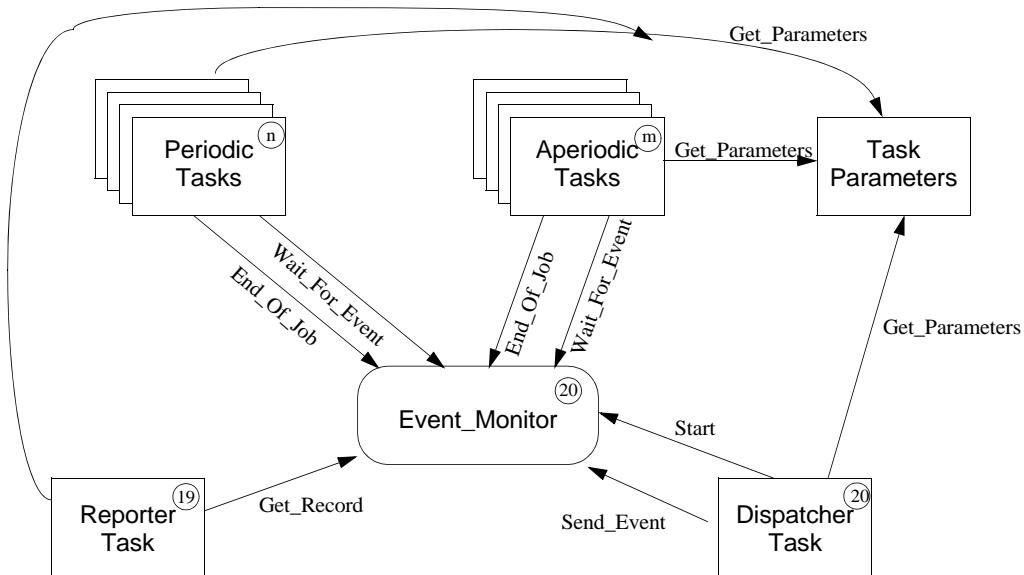
## 2. Periodic tasks

- Assignment 2.1.a: Periodic Tasks in Ada
- Assignment 2.1.b: Periodic Tasks in C/POSIX

## Assignment 2.1.a Periodic Tasks in Ada

1. Develop an application with four periodic tasks, with utilization less than U(4), using a dispatcher task
  - create a main program with the four tasks
  - define the parameters in `task_params.adb`
2. Assign RMA priorities and observe the response times
3. Assign different priorities and check if the deadlines are met or not
4. With the RMA priorities, increase the utilization until the deadlines are not met
5. Write a report with:
  - source code of the tasks and main program
  - observations on response times under the different cases

# Software Architecture



# Software Architecture

## Three tasks:

- **User tasks:** periodics (and later, aperiodics)
- **Dispatcher:** Generates the events that cause the execution of each user task job
- **Reporter:** generates a report with the results of execution

## One protected object

- **Event monitor:** Stores the events and task job information

# Task Parameters

```
with System;
package Task_Params is

    type Task_Count is range 0..10;
    subtype Task_ID is Task_Count range 1..10;

    type Task_Parameters is record
        Period, Exec_Time, Deadline : Duration;
        ... -- other parameters
    end record;

    function Get_Parameters(Of_Task : Task_Id)
        return Task_Parameters;

    function Number_Of_Tasks return Task_Count;

end Task_Params;
```

# Dispatcher Task

```
task Dispatcher is
    pragma Priority (20);
end Dispatcher;
```

# Dispatcher Task (cont'd)

```

task body Dispatcher is
  Activation_Queue  -- store task ids
    -- and activation times
  Initial_Time : Time:=Clock;
begin
  for The_Task in All_Tasks loop
    Activation_Queue.Enqueue(The_Task,Initial_Phase);
  end loop;
  Monitor.Start(Initial_Time);
  loop
    Activation_Queue.Dequeue(The_Task,Next_Activation);
    delay until (Initial_Time+Next_Activation);
    Monitor.Send_Event(The_Task);
    Activation_Queue.Enqueue(The_Task,
      Next_Activation+Period(The_Task)); -- for periodics
  end loop;
end Dispatcher;

```

# Execution Time Load

```

package Execution_Load is

  procedure Eat (For_Seconds : Duration);

end Execution_Load;

```

# Event Monitor

---

```

with Ada.Calendar, Task_Params;
package Event_Monitor is
  ...
  protected Monitor is
    procedure Start (Initial_Time : in Time);
    procedure Send_Event (The_Task : in Task_ID);
    entry Wait_For_Event (Task_ID);--family of entries
    procedure End_Of_Job (The_Task : in Task_ID);
    procedure Put_Message (Msg   : in String);
    ...
    pragma Priority (20);
  private
    ...
  end Monitor;
end Event_Monitor;

```

# Structure of a Periodic Task

---

```

task type Periodic
  (ID : Task_Id;
   Prio : Priority)
is
  pragma Priority(Prio); -- must be less than 19
end Periodic;

task body Periodic is
begin
  loop
    Monitor.Wait_For_Event(ID);
    Execution_Load.Eat(Exec_Time);
    Monitor.End_Of_Job(ID);
  end loop;
end Periodic;

```

# Assignment 2.1.b: Periodic Tasks in C/POSIX

1. Develop an application with four periodic threads, with utilization less than U(4)
2. Assign RMA priorities and observe the response times
3. Assign different priorities and check if the deadlines are met or not
4. With the RMA priorities, increase the utilization until the deadlines are not met
5. Write a report with
  - observations on response times under the different cases

## Execution Time Load

### File eat.h:

```
void eat(const struct timespec *cpu_time);
```

### POSIX struct timespec

```
struct timespec {  
    time_t tv_sec; // seconds  
    long tv_nsec; // nanoseconds  
}
```

- measures absolute time since the Epoch,
- or a relative interval

# Main Program

---

The main program does:

- set its own priority to a very high value
  - so that the other threads start at the same time, when the main goes to sleep
- create a threads attributes object and set:
  - inheritsched attribute to **PTHREAD\_EXPLICIT\_SCHED**
  - detach state to **PTHREAD\_CREATE\_DETACHED**
  - scheduling policy to **SCHED\_FIFO**
- create four periodics threads; for each:
  - change the priority
  - create the thread
- sleep for a very long time

# Structure of a Periodic Thread

---

```
periodic_thread (id, period, exec_time) {
    activation_time = read the clock
    while (1) {
        report ("START" + id + current time)
        eat (exec_time)
        response_time=clock_gettime-activation_time
        report ("END" + id + response_time)
        activation_time=activation_time+period
        sleep until (activation_time)
    }
}
```

### 3. Extensions to the periodic model

- Assignment 3.1: Simulating an interrupt

## Assignment 3.1: Simulating an interrupt

Modify the implementation of periodic tasks to simulate an interrupt service routine

- By changing a medium priority task to the highest priority level

Compare the response times with those obtained when rate monotonic priorities were used

- what conclusions can you derive?

Now try simulating the following solution

- leave part of the work at interrupt priority
- move another part of the work to a new task with rate monotonic priority
- compare and discuss the response times obtained

## 4. Modelling real-time systems

- Assignment 4.1: Model and analyze the system developed in Assignment 2.1 or 2.2
- Assignment 4.2: Model and analyze the system developed in Assignment 3.1

## Assignment 4.1: Model and analyze the system of Assignment 2.1

Install the MAST toolset and graphical editor

Create a model to represent the system developed in Assignment 2.1 or 2.2

Analyze the response times and check them against the previous results

Write a short report with the results of the analysis

# Assignment 4.2: Model and analyze the system of Assignment 3.1



Create a model to represent the system developed in Assignment 3.1

Analyze the response times and check them against the previous results

Write a short report with the results of the analysis



## 5. Synchronization

- Assignment 5.1.a: Implement task synchronization in the periodic task simulator in Ada
- Assignment 5.1.b: Implement task synchronization in the periodic task simulator in C/POSIX
- Assignment 5.2: Model and analyze periodic tasks with shared objects

## Assignment 5.1.a: Synchronization

In the task simulation application, create two protected objects to control access to two shared objects

- modify three of the tasks to use these protected objects (one of them using both)
- leave a higher priority task that does not use resources

Do not assign priority ceilings

- the default priority ceiling in Ada is the highest priority
- play with the task phases to achieve the highest blocking times

Assign the appropriate priority ceilings and check the response times

Write a report containing the code of the tasks, protected objects, and main program. Also add a short summary of the results.

# Shared Resource Parameters

---

```

with System;
package Resource_Params is

  type Resource_Count is range 0..2;
  subtype Resource_Id is Resource_Count
    range 1..Resource_Count'Last;

  type Section_Id is range 1..2;

  type Exec_Times is array (Section_Id) of Duration;

  type Resource_Parameters is record
    Prio : System.Priority;
    Exec : Exec_Times;
  end record;

```

# Shared Resource Parameters (cont'd)

---

```

function Get_Parameters
  (Of_Resource : Resource_Id)
  return Resource_Parameters;

function Number_Of_Resources
  return Resource_Count;

end Resource_Params;

```

# Specification of the Protected Object

```

protected type Shared_Resource(Ceiling :Priority) is

  procedure Critical_Sect_1 (Length : in Duration);
  procedure Critical_Sect_2 (Length : in Duration);

  pragma Priority (Ceiling);
  // Comment out the line above to use default ceiling

end Shared_Resource;

```

# Body of the Protected Object

```

protected body Shared_Resource is

  procedure Critical_Sect_1
    (Length : in Duration) is
  begin
    Execution_Load.Eat(Length);
  end Critical_Sect_1

  procedure Critical_Sect_2
    (Length : in Duration) is
  begin
    Execution_Load.Eat(Length);
  end Critical_Sect_2

end Shared_Resource;

```

# Periodic Task that Uses Shared Objects

```
task body Periodic_With_Synch is
begin
    Exec_Time:=Exec_Time-Length of critical section
    loop
        Monitor.Wait_For_Event(ID);
        Execution_Load.Eat(Exec_Time/2);

        Resource_Critical_Sect_1(Length);

        Execution_Load.Eat(Exec_Time/2);
        Monitor.End_Of_Job(ID);
    end loop;
end Periodic;
```

## Assignment 5.1.b: Implement task synchronization in C/POSIX

A task simulation application that uses one mutex to control access to a shared object

- the shared object is coded in a monitor with operations that use the mutex
- the mutex uses the immediate priority ceiling protocol
  - the ceiling is adjustable at initialization time
- two tasks use operations of the shared object

## Assignment 5.1.b (cont'd)

Adapt this application to experiment with priority inversion

- create a medium priority task that does not use the shared object
- change the mutex protocol to PTHREAD\_PRIO\_NONE
- play with the task phases to observe the unbounded priority inversion effect

Then

- use the PTHREAD\_PRIO\_PROTECT protocol for the mutex
- assign the appropriate priority ceilings
- check the response times

Write a report with the developed code and a short summary of the results

## Assignment 5.2: Model and analyze periodic tasks with shared objects

Build a MAST model of the simulation application with shared resources, and analyze it

Compare the results with those observed in practice

Write a report with the results obtained, including the response times and blocking times

## 6. Aperiodic events

- Assignment 6.1.a. Aperiodic tasks in Ada
- Assignment 6.1.b. Aperiodic tasks in C/POSIX

## Assignment 6.1.a: Aperiodic tasks in Ada

- Create one aperiodic task
- Observe the unpredictability introduced by this task
- Use the polling server approach, and check the predictability of the response times
- Use the emulated sporadic server approach
- Build a MAST model of the system and obtain the response times

# Specification of aperiodic tasks

```

with System;
with Resource_Params;

package Task_Params is

  type Task_Count is range 0..10;
  subtype Task_Id is Task_Count range 1..Task_Count'Last;

  type Event_Kinds is (Periodic, Aperiodic);

  type Task_Kinds is (Regular, With_Sync, Polling, Sporadic);

```

# Specification of aperiodic tasks (cont'd)

```

type Task_Parameters is record
  Event_Kind : Event_Kinds;
  Task_Kind : Task_Kinds;
  Prio : System.Priority;
  Period, Exec_Time, Deadline, Phase,
    Avg_Interarrival, Initial_Capacity : Duration;
  Res_Id : Resource_Params.Resource_ID;
  Sect_Id : Resource_Params.Section_Id;
end record;

function Get_Parameters (Of_Task : Task_Id)
  return Task_Parameters;

function Number_Of_Tasks return Task_Count;

end Task_Params;

```

# Simple Aperiodic Task

```

task type Aperiodic
  (ID    : task_id;
   Prio  : Priority)
is
  pragma Priority(Prio);
end Aperiodic;

task body Aperiodic is
begin
  loop
    Monitor.Wait_For_Event(ID);
    Execution_Load.Eat(Exec_Time);
    Monitor.End_Of_Job(ID);
  end loop;
end Aperiodic;

```

# Aperiodic Task with Polling

```

task body Aperiodic is
  Next_Activation : Time:=Clock;
begin
  loop
    select
      Monitor.Wait_For_Event(ID);
      Execution_Load.Eat(Exec_Time);
      Monitor.End_Of_Job(ID);
    else
      null;
    end select;
    Next_Activation:=Next_Activation+Period;
    delay until Next_Activation;
  end loop;
end Aperiodic;

```

# Aperiodic Task with Emulated Sporadic Server

```
task body Aperiodic is
    Capacity : Duration:=Initial_Capacity;
begin
    loop
        Monitor.Wait_For_Event(ID);
        Last_Activation:=Clock;
        Execution_Load.Eat(Exec_Time);
        Monitor.End_Of_Job(ID);
        Capacity:=Capacity-Exec_Time;
        if Capacity<Exec_Time then
            delay until Last_Activation+Period;
            Capacity:=Initial_Capacity;
        end if;
    end loop;
end Aperiodic;
```

## Assignment 6.1.b: Aperiodic tasks in C/POSIX

- Use the proposed application with one aperiodic thread and two periodic threads
  - the events arrive with a random interarrival time
- Observe the unpredictability introduced by this thread
- Use the polling server approach, and check the predictability of the response times
- Use the emulated sporadic server approach
- Build a MAST model of the system and obtain the response times

# Structure of an Aperiodic Thread

```

aperiodic_thread (id, signal, avg_interarrival,
                  exec_time)
{
    next_event= initial_time+phase;
    while (1) {
        // simulate the wait for an aperiodic event
        clock_nanosleep(ABS_TIME,next_event);
        // process the event
        report ("START" + id + current period)
        eat (exec_time)
        response_time=clock_gettime-next_event
        report ("END" + id + response_time)
        // simulate arrival of next event
        next_event= next_event+
                    random(0..2*avg_interarrival);
    }
}
  
```

# Structure of an Aperiodic Thread with Polling

```

aperiodic_thread (id, signal, avg_interarrival, polling_period,
                  exec_time)
{
    next_event= initial_time+phase;
    next_activation= next_event
    next_event= next_event+random(0..2*avg_interarrival);
    while (1) {
        // simulate the arrival of an aperiodic event
        now=clock_gettime
        if (now>=next_event) {
            // event has arrived
            report ("START" + id + current period)
            eat (exec_time)
            response_time=clock_gettime-next_event
            report ("END" + id + response_time)
            next_event=next_event+random(0..2*avg_interarrival);
        }
    }
}
  
```

# Structure of an Aperiodic Thread with Polling

```
// wait for next period
next_activation=next_activation+polling_period
clock_nanosleep(ABS_TIME,next_activation);
}
}
```

# Structure of an Aperiodic Thread with Emulated Sporadic Server

```
aperiodic_thread (id, signal, avg_interarrival,
                  initial_capacity, repl_period, exec_time)
{
    next_event= initial_time+phase;
    capacity=initial_capacity
    while (1) {
        // simulate the wait for an aperiodic event
        next_event= next_event+
            random(0..2*avg_interarrival);
        clock_nanosleep(ABS_TIME,next_event);
        last_activation=clock_gettime();
        // process the event
        report ("START" + id + current period)
        eat (exec_time)
        response_time=clock_gettime-next_event
        report ("END" + id + response_time)
        capacity=capacity-exec_time
    }
}
```

# Structure of an Aperiodic Thread with Emulated Sporadic Server

```
// check available capacity
if (capacity<exec_time) {
    clock_nanosleep(ABS_TIME,
                    last_Activation+repl_period)
    capacity=initial_capacity
}
}
```

# Modelling and Analysis of Real-Time Systems: Lab Assignments



## 7. Dynamic priority scheduling

- Assignment 7. Use MAST to model and analyze the system described in Exercise 7



## 8. Advanced topics

- Assignment 8. Model and analyze a system with periodic tasks with release jitter and post-period deadlines.

## Assignment 8. Release Jitter and Post-Period Deadlines

Suppose the following system:

	C	T	D	J
Task $\tau_1$ :	20	100	100	10
Task $\tau_2$ :	25	150	150	10
Task $\tau_3$ :	100	350	700	10

We now want to add a new periodic transaction that does the following actions

- RS: read a sensor
- PR: process the information from the sensor
- OA: output the result to an actuator
- The transaction is periodic, with a period of 50 ms and 1ms release jitter

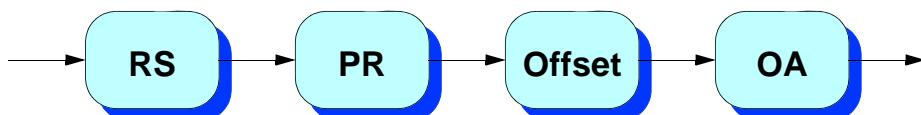
## Assignment 8 (cont'd)

The requirements of the new transaction are

	C	D	Max Output Jitter
RS	1	-	3
PR	5	-	
OA	1	20	3

## Assignment 8 (cont'd)

To accomplish these requirements we design the transaction as follows



The offset is set to the worst-case response time of activity PR.

RS and OA are given the highest priority

PR is given the next highest priority

# Assignment 8 (cont'd)

## Steps:

- create a MAST model of this system
- analyze it to obtain the correct value of the offset
- make a final analysis to check whether the system is schedulable or not
- write a short report with the results obtained

# Appendix 1: Using the Gnat Compiler

To compile with gnat-2007:

```
• /usr/local/gnat-2007/gnatvars.sh
```

To compile a module called **name.adb**:

```
gcc -c -gnatv name.adb
```

To compile everything and link the program called **name**

```
gnatmake -gnatv name.adb
```

The same, using the MaRTE OS run-time system

```
gnatmake -gnatv --RTS=marte name.adb
```

# Appendix 1: Using the Gnat Compiler

Main compiler options

- **-gnatv**: shows the lines with errors
- **-g**: needed to use the debugger
- **-gnato**: to check the results of arithmetic operations
- **example**: **gnatmake -gnatv -gnato -g name.adb**

## Appendix 2: Using MaRTE OS in Linux

Set the environment to compile with **MaRTE OS** under the **linux** architecture:

- `/usr/local/marte-1.8/martevars.sh`

Compile a C module called **name.c**:

```
mgcc -c name.c
```

Compile a C program called **prog.c** and link it with a module called **name.o**:

```
mgcc prog.c name.o
```

Execute the program:

```
./mprogram
```

## Appendix 3: Using MAST

To have access to MAST:

- `/usr/local/mast-1.4.0.1/mastvars.sh`